

REALTINE OPERATING SYSTEMS

REALTIME OPERATING System is said to be Real Time if it is required to complete it's work & deliver it's services on time. Example – Flight Control System All tasks in that system must execute on time. Non Example – PC system

HARDAND SOFT REAL Hard Refer TSNE Systems

Validation by provably correct procedures or extensive simulation that the system always meets the timings constraints Failure to meet deadlines is fatal example : Flight Control System

SOFT REALTIME SYSTEM

- Demonstration of jobs meeting some statistical constraints suffices.
 - Late completion of jobs is undesirable but not fatal.
 - System performance degrades as more & more jobs miss deadlines
 - Online Databases

ROLE OF AN OS IN BEAL THE STENS Often no OS involved Micro controller based Embedded Systems

Some Real Time Applications are huge & complex

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^DMultiple threads

Complicated Synchronization Requirements Filesystem / Network / Windowing support OS primitives reduce the software design time



* Resource Allocation.

Interrupt Handling.

Other issues like kernel size.

<u>SCHEDULING IN</u> <u>RIOS</u>

More information about the tasks are known No of tasks Resource Requirements Release Time Execution time Deadlines

Being a more deterministic system better scheduling algorithms can be devised.

SCHEDULING ALGARICHIMS IN RTOS Weighted Round Robin Scheduling

Priority Scheduling (Greedy / List / Event Driven) CLOCK DRVEN All parameters about jobs (release time/ execution time/deadline) known in advance. Schedule can be computed offline or at some regular time instances. Minimal runtime overhead. Not suitable for many applications.

MEIGHTED ROUND ROBIN

- Jobs scheduled in FIFO manner
- Time quantum given to jobs is proportional to it's weight
- Example use : High speed switching network QOS guarantee.
- Not suitable for precedence constrained jobs. Job A can run only after Job B. No point in giving time quantum to Job B before Job A.

PRIORITY SCHEDULING

(Greedy/List/Event Driven) Processor never left idle when there are ready tasks Processor allocated to processes according to priorities Priorities

static - at design time

Dynamic - at runtime

PRORTY Easiest beachelerst Ceder) Process with earliest deadline given highest priority Least Slack Time First (LSF) slack = relative deadline - execution left Rate Monotonic Scheduling (RMS) For periodic tasks Tasks priority inversely proportional to it's period

RESOURCE RALL-ORIATION IN The issues with scheduling applicable here. Resources can be allocated in Weighted Round Robin Priority Based

Some resources are non preemptible Example : semaphores Priority Inversion if priority scheduling is used

INTERRUPT Internel and the same an isocritical sections of the kernel No worst case bound on interrupt latency available g: Disk Drivers may disable interrupt for few hundred milliseconds Not suitable for Real Time Applications Interrupts may be missed

TWOLEVELNDERRUPT TRANSPICE Handling Top Half Interrupt Handler Called Immediately – Kernel never disables interrupts Cannot invoke thread library functions - Race Conditions Bottom Half Interrupt Handler Invoked when kernel not in Critical Section Can invoke thread library functions Very Low Response time (as compared to Linux)

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<u>OTHERFEATURES</u>

Footprint Small footprint (~50kb)

Oskit's Device Driver Framework Allows direct porting of existing drivers from Linux. Example – Ethernet Driver of Linux